# Connectivity Definitions and Theorems from Weeks 3 and 4 

## February 13

Connectivity Definitions:
Definition: A graph $G$ is connected if for all $v, w \in V(G)$, there exists a path from $v$ to $w$.
Definition: $G$ is disconnected if $G$ is not connected.
Definition: A (connected) component $H$ is a maximal subgraph $H$ of $G$ that is connected.
Definition: $G$ is $\underline{k-c o n n e c t e d ~ i f ~}|V(G)|>k$ and removing fewer than $k$ vertices does not disconnect the graph.
(We say that every graph is 0 -connected.)
Definition: The connectivity of $G$ (denoted $\kappa(G)=$ "kappa") is the maximum $k$ such that $G$ is $k$-connected.
(Conventions: The connectivity of a single vertex is zero and $\kappa\left(K_{n}\right)=n-1$.)
Definition: A cut vertex is a vertex $v \in V(G)$ such that $G \backslash v$ is disconnected.
Definition: A separating set or vertex cut is a set of vertices $X \subset V(G)$ such that $G \backslash X$ is disconnected.
Note: $\kappa(G)=0 \Longleftrightarrow G$ is disconnected or $G$ is a single vertex
Note: $\kappa(G) \geq 2 \Longleftrightarrow G$ has no cut vertex.
Edge Connectivity Definitions:
Definition: $G$ is $k$-edge-connected if removing fewer than $k$ edges does not disconnect the graph.
(We say that every graph is 0-edge-connected.)
Definition: The edge connectivity of $G$ (denoted $\lambda(G)=$ "lambda" or $\kappa^{\prime}(G)$ ) is the maximum $k$ such that $G$ is $k$-edge connected.
Definition: A bridge is an edge $e \in E(G)$ such that $G \backslash e$ is disconnected.
Definition: A DISCONNECTING SET is a set of edges $D \subset E(G)$ such that $G \backslash D$ is disconnected.
Note: $\lambda(G)=0 \Longleftrightarrow G$ is disconnected.
Note: $\lambda(G) \geq 2 \Longleftrightarrow G$ has no bridge.

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Note: If you delete a cut vertex from a graph, the number of connected components increases.
Note: If you delete a bridge from a graph, the number of connected components increases by exactly one.
Theorems:
Theorem: (Book 2.4.1) Let $G$ be connected. Then $G$ is a tree $\Longleftrightarrow$ Every edge of $G$ is a bridge.
Theorem: (Book 3.2.1) A regular graph of even degree has no bridge.
Theorem: For all graphs $G, \lambda(G) \leq \delta(G)$.
Edge Cuts: [not to be confused with cutset or cut vertex]
Definition: Let $X \subset V(G)$. Then $X^{c}$ is the complement of $X$.
That is, $V(G)=X \cup X^{c}$ and $X \cap X^{c}=\emptyset$.
Definition: For any $X \subset V(G)$ such that $X, X^{c} \neq \emptyset$, an edge cut (denoted $\left[X, X^{c}\right]$ ) is the set of edges $D$ between $X$ and $X^{c}$.
Note: An edge cut is a disconnecting set, but not vice versa.
(This implies $\lambda(G) \leq\left|\left[X, X^{c}\right]\right|$ for all $X \subset V(G)$.
Note: A minimal disconnecting set is an edge cut, but not vice versa.
Theorem: For all graphs $G, \kappa(G) \leq \lambda(G)$.

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## Blocks:

Definition: A block of a graph $G$ is a maximally connected subgraph of $G$ with no cut vertex.
Note: The following things are true about blocks.

1. $G$ itself may be a block.
2. Except for blocks that are edges, blocks are always 2-connected.
3. Any two blocks share at most one vertex.
4. A vertex shared between blocks is a cut vertex of $G$.
5. The blocks of $G$ partition $E(G)$.

Definition: The block graph of $G$ is a bipartite graph $H$ with vertices $v_{i}$ representing cut vertices of $G$, and vertices $b_{j}$ representing blocks of $G$, where $v_{i} b_{j}$ is an edge of $H$ if vertex $v_{i}$ is a vertex in block $b_{j}$.
Note: A block graph is always a forest. (Proof in hwk.)
More Graph Statistics:
Definition: An independent set of a graph $G$ is a subset $X \subset V(G)$ such that no edge of $G$ connects any two vertices of $X$. In other words, the induced subgraph of $G$ on $X$ contains no edges.
Definition: The independence number of a graph $G$ is the size of the maximum independence set of $G$. It is denoted $\alpha(G)$.
Definition: A vertex cover of a graph $G$ is a subset $X \subset V(G)$ such that $X$ contains (at least) one endpoint of every edge in $G$.
Definition: The size of the smallest vertex cover is denoted $\beta(G)$.
Theorem: In any graph $G, X \subset V(G)$ is an independent set $\Longleftrightarrow X^{c}$ is a vertex cover.
Theorem: For all graphs $G, \alpha(G)+\beta(G)=|V(G)|$.
Note: Finding an independent set $X$ and a vertex cover $Y$ such that $|X|+|Y|=|V(G)|$ implies that $\alpha(G)=|X|$ and $\beta(G)=|Y|$.
Definition: The clique number $\omega(G)$ of a graph $G$ is the largest number $k$ such that $K_{k}$ is a subgraph of $G$.

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Characterization of 2-connectedness:
Theorem: (Whitney, 1932 aka MINI-Menger) Let $G$ be a graph with $\geq 3$ vertices. Then $G$ is 2-connected $\Longleftrightarrow$ for all $v, w \in V(G)$, there exist two internally disjoint $v, w$-paths in $G$.
Theorem: (Menger) $G$ is $k$-connected $\Longleftrightarrow$ for all $v, w \in V(G)$, there exist $k$ internally disjoint $v, w$-paths in $G$.
 and ends in $H$.
Definition: An ear decomposition is a construction of $G$ starting with some cycle $C$, and at each step successively adding to the existing graph $H$ some $H$-path.
Theorem: $G$ is 2-connected $\Longleftrightarrow G$ has an ear decomposition.
Theorem: Let $G$ be a graph with $\geq 3$ vertices. The following are equivalent:

1. $G$ is 2-connected.
2. $G$ is connected and has no cut vertex.
3. $G$ is a block.
4. For all $v, w \in V(G)$, there exist two internally disjoint $v, w$-paths in $G$.
5. For all $v, w \in V(G)$, there exists a cycle in $G$ through $v$ and $w$.
6. $\delta(G)>0$ and for all $e, f \in E(G)$, there exists a cycle in $G$ through $e$ and $f$.
7. $G$ has an ear decomposition.

## List of graph statistics so far

$\alpha(G)$
$\beta(G)$
$\delta(G)$
$\Delta(G)$
$\kappa(G)$
$\lambda(G)$
$\omega(G)$
$\operatorname{diam}(G)$
$g(G)$

