# Connectivity Definitions and Theorems from Weeks 3 and 4

## February 13

Connectivity Definitions:

**Definition:** A graph G is <u>connected</u> if for all  $v, w \in V(G)$ , there exists a path from v to w.

**Definition:** G is disconnected if G is not connected.

**Definition:** A (connected) component H is a MAXIMAL subgraph H of G that is connected.

**Definition:** G is  $\underline{k\text{-connected}}$  if |V(G)| > k and removing fewer than k vertices does not disconnect the graph.

(We say that every graph is 0-connected.)

**Definition:** The <u>connectivity</u> of G (denoted  $\kappa(G) =$  "kappa") is the maximum k such that G is k-connected.

(Conventions: The connectivity of a single vertex is zero and  $\kappa(K_n) = n - 1$ .)

**Definition:** A cut vertex is a vertex  $v \in V(G)$  such that  $G \setminus v$  is disconnected.

**Definition:** A <u>SEPARATING SET</u> or <u>vertex cut</u> is a set of vertices  $X \subset V(G)$  such that  $G \setminus X$  is disconnected.

Note:  $\kappa(G) = 0 \iff G$  is disconnected or G is a single vertex

**Note:**  $\kappa(G) \geq 2 \iff G$  has no cut vertex.

Edge Connectivity Definitions:

**Definition:** G is  $\underline{k}$ -edge-connected if removing fewer than k edges does not disconnect the graph.

(We say that every graph is 0-edge-connected.)

**Definition:** The edge connectivity of G (denoted  $\lambda(G)$  = "lambda" or  $\kappa'(G)$ ) is the maximum k such that G is k-edge connected.

**Definition:** A bridge is an edge  $e \in E(G)$  such that  $G \setminus e$  is disconnected.

**Definition:** A DISCONNECTING SET is a set of edges  $D \subset E(G)$  such that  $G \setminus D$  is disconnected.

Note:  $\lambda(G) = 0 \iff G$  is disconnected. Note:  $\lambda(G) > 2 \iff G$  has no bridge.

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Note: If you delete a cut vertex from a graph, the number of connected components increases.

**Note:** If you delete a bridge from a graph, the number of connected components increases by exactly one.

#### Theorems:

**Theorem:** (Book 2.4.1) Let G be connected. Then G is a tree  $\iff$  Every edge of G is a bridge.

**Theorem:** (Book 3.2.1) A regular graph of even degree has no bridge.

**Theorem:** For all graphs G,  $\lambda(G) \leq \delta(G)$ .

Edge Cuts: [not to be confused with cutset or cut vertex]

**Definition:** Let  $X \subset V(G)$ . Then  $X^c$  is the complement of X.

That is,  $V(G) = X \cup X^c$  and  $X \cap X^c = \emptyset$ .

**Definition:** For any  $X \subset V(G)$  such that  $X, X^c \neq \emptyset$ , an edge cut (denoted  $[X, X^c]$ ) is the set of edges D between X and  $X^c$ .

Note: An edge cut is a disconnecting set, but not vice versa.

(This implies  $\lambda(G) \leq |[X, X^c]|$  for all  $X \subset V(G)$ .

Note: A minimal disconnecting set is an edge cut, but not vice versa.

**Theorem:** For all graphs G,  $\kappa(G) \leq \lambda(G)$ .

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Blocks:

**Definition:** A <u>block</u> of a graph G is a maximally connected subgraph of G with no cut vertex. **Note:** The following things are true about blocks.

- 1. G itself may be a block.
- 2. Except for blocks that are edges, blocks are always 2-connected.
- 3. Any two blocks share at most one vertex.
- 4. A vertex shared between blocks is a cut vertex of G.
- 5. The blocks of G partition E(G).

**Definition:** The block graph of G is a bipartite graph H with vertices  $v_i$  representing cut vertices of G, and vertices  $\overline{b_j}$  representing blocks of G, where  $v_ib_j$  is an edge of H if vertex  $v_i$  is a vertex in block  $b_j$ .

**Note:** A block graph is always a forest. (Proof in hwk.)

More Graph Statistics:

**Definition:** An independent set of a graph G is a subset  $X \subset V(G)$  such that no edge of G connects any two vertices of X. In other words, the induced subgraph of G on X contains no edges.

**Definition:** The independence number of a graph G is the size of the maximum independence set of G. It is denoted  $\alpha(G)$ .

**Definition:** A <u>vertex cover</u> of a graph G is a subset  $X \subset V(G)$  such that X contains (at least) one endpoint of every edge in G.

**Definition:** The size of the smallest vertex cover is denoted  $\beta(G)$ .

**Theorem:** In any graph  $G, X \subset V(G)$  is an independent set  $\iff X^c$  is a vertex cover.

**Theorem:** For all graphs G,  $\alpha(G) + \beta(G) = |V(G)|$ .

**Note:** Finding an independent set X and a vertex cover Y such that |X| + |Y| = |V(G)| implies that  $\alpha(G) = |X|$  and  $\beta(G) = |Y|$ .

**Definition:** The clique number  $\omega(G)$  of a graph G is the largest number k such that  $K_k$  is a subgraph of G.

# February 20

Characterization of 2-connectedness:

**Theorem:** (Whitney, 1932 aka MINI-Menger) Let G be a graph with  $\geq 3$  vertices. Then G is 2-connected  $\iff$  for all  $v, w \in V(G)$ , there exist two internally disjoint v, w-paths in G.

**Theorem:** (Menger) G is k-connected  $\iff$  for all  $v, w \in V(G)$ , there exist k internally disjoint v, w-paths in G.

**Definition:** Let H be any subgraph of G. Then an  $\underline{H\text{-path}}$  (or an  $\underline{\operatorname{ear}}$ ) is a path in G that starts and ends in H.

**Definition:** An <u>ear decomposition</u> is a construction of G starting with some cycle C, and at each step successively adding to the existing graph H some H-path.

**Theorem:** G is 2-connected  $\iff$  G has an ear decomposition.

**Theorem:** Let G be a graph with  $\geq 3$  vertices. The following are equivalent:

- 1. G is 2-connected.
- 2. G is connected and has no cut vertex.
- 3. G is a block.
- 4. For all  $v, w \in V(G)$ , there exist two internally disjoint v, w-paths in G.
- 5. For all  $v, w \in V(G)$ , there exists a cycle in G through v and w.
- 6.  $\delta(G) > 0$  and for all  $e, f \in E(G)$ , there exists a cycle in G through e and f.
- 7. G has an ear decomposition.

# List of graph statistics so far

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