## Bipartite graphs

Question: What is $\chi\left(C_{n}\right)$ when $n$ is odd?
Answer:
Definition: A graph is called bipartite if $\chi(G) \leq 2$.
Examples: $K_{m, n}, \square_{n}$, Trees

Theorem 2.1.6: $G$ is bipartite $\Longleftrightarrow$ every cycle in $G$ has even length.
$(\Rightarrow)$ Let $G$ be bipartite. Assume that there is some cycle $C$ of odd length contained in G...

## Proof of Theorem 2.1.6

$(\Leftarrow)$ Suppose that every cycle in $G$ has even length. We want to show that $G$ is bipartite. Consider the case when $G$ is connected.

Plan: Construct a coloring on $G$ and prove that it is proper.
Choose some starting vertex $x$ and color it blue. For every other vertex $y$, calculate the distance from $y$ to $x$ and then color $y$ :

$$
\begin{cases}\text { blue } & \text { if } d(x, y) \text { is even. } \\ \text { red } & \text { if } d(x, y) \text { is odd. }\end{cases}
$$

Question: Is this a proper coloring of $G$ ?
Suppose not. Then there are two vertices $v$ and $w$ of the same color that are adjacent. This generates a contradiction because there exists an odd cycle as follows:

## Edge Coloring

Parallel to the idea of vertex coloring is the idea of edge coloring.
Definition: An edge coloring of a graph $G$ is a labeling of the edges of $G$ with colors. [Technically, it is a function $f: E(G) \rightarrow\{1,2, \ldots, /\}$.]

Definition: A proper edge coloring of $G$ is an edge coloring of $G$ such that no two adjacent edges are colored the same.

Example: Cube graph $\left(\square_{3}\right)$ :


We can properly edge color $\square_{3}$ with $\qquad$ colors and no fewer.

Definition: The minimum number of colors necessary to properly edge color a graph $G$ is called the edge chromatic number of $G$, denoted $\chi^{\prime}(G)=$ "chi prime".

## Edge coloring theorems

Theorem 2.2.1: For any graph $G, \chi^{\prime}(G) \geq \Delta(G)$.
Theorem 2.2.2: Vizing's Theorem:
For any graph $G, \chi^{\prime}(G)$ equals either $\Delta(G)$ or $\Delta(G)+1$.
Proof: Hard. (See reference [24] if interested.)
Consequence: To determine $\chi^{\prime}(G)$,

Fact: Most 3-regular graphs have edge chromatic number 3.


## Snarks

Definition: A 3-regular graph with edge chromatic number 4 is called a snark.

Example: The Petersen graph $P$ :


The edge chromatic number of complete graphs
Goal: Determine $\chi^{\prime}\left(K_{n}\right)$ for all $n$.
Vertex Degree Analysis: The degree of every vertex in $K_{n}$ is $\qquad$ .

Vizing's theorem implies that $\chi^{\prime}\left(K_{n}\right)=$ $\qquad$ or $\qquad$ .
If $\chi^{\prime}\left(K_{n}\right)=\ldots$, then each vertex has an edge leaving of each color.
Q: How many red edges are there?
This is only an integer when:
So, the best we can expect is that $\left\{\begin{array}{l}\chi^{\prime}\left(K_{2 n}\right)= \\ \chi^{\prime}\left(K_{2 n-1}\right)=\end{array}\right.$

The edge chromatic number of complete graphs
Theorem 2.2.3: $\quad \chi^{\prime}\left(K_{2 n}\right)=2 n-1$.
Proof: We prove this using the turning trick.
Label the vertices of $K_{2 n}$
$0,1, \ldots, 2 n-2, x$. Now,
Connect 0 with $x$
Connect 1 with $2 n-2$,
Connect $n-1$ with $n$.
Now turn the edges.
And do it again. (and again, ...)
Each time, new edges are used.
This is because each of the
 edges is a different "circular length": vertices are at circ. distance $1,3,5, \ldots, 4,2$ from each other, and $x$ is connected to a different vertex each time.

The edge chromatic number of complete graphs
Theorem 2.2.4: $\quad \chi^{\prime}\left(K_{2 n-1}\right)=2 n-1$.
This construction also gives a way to edge color $K_{2 n-1}$ with $2 n-1$ colors-simply delete vertex $x$ !

This is related to the area of combinatorial designs.
Question: Is it possible for six tennis players to play one match per day in a five-day tournament in such a way that each player plays each other player once?
Day $1 \quad 0 x \quad 14 \quad 23$

Day 2 1x 2034
Day 3 2x 3140
Day 4 3x $42 \quad 01$
Day 5 4x 0312


Theorem 2.2.3 proves there is such a tournament for all even numbers.

