

# The crossing number of a graph

Some graphs are *almost* planar.

- ▶ If  $K_{3,3}$  didn't have that last edge, it would be planar!

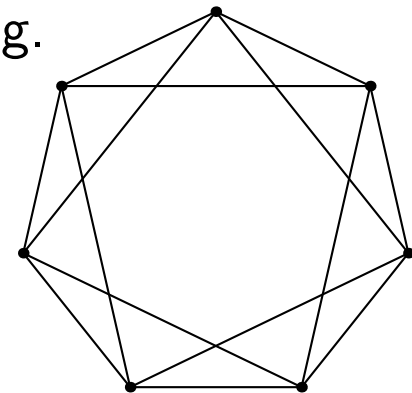
**So we ask:** How non-planar is it?

★ We will discuss **three** ways to answer this question.

*Definition:* The **crossing number** of a graph  $G$ , denoted  $\text{cr}(G)$ , is the minimum number of crossings in any simple drawing of  $G$ .

- ▶ So if  $G$  is planar,  $\text{cr}(G) = 0$ , and if  $G$  is non-planar,  $\text{cr}(G) \geq 1$ .
- ▶ To prove  $\text{cr}(G) = 1$ :
  - ▶ Prove  $G$  is non-planar (Kuratowski or otherwise) **and**
  - ▶ Find a drawing of  $G$  with only one crossing.

**Example.**

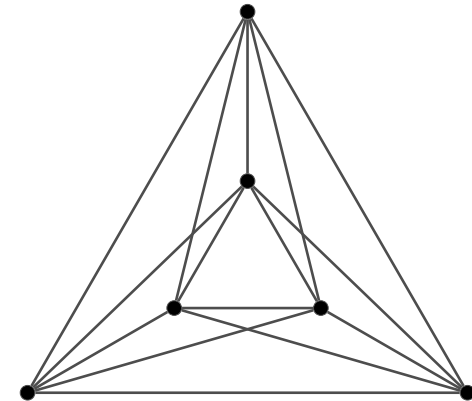


# The crossing number of $K_6$

*Theorem 9.1.4* The crossing number of  $K_6$  is 3.

*Proof.* First, here is a drawing of  $K_6$  with three crossings:

We conclude that  $\text{cr}(K_6) \leq 3$ .



*Claim:* No simple drawing of  $K_6$  has fewer crossings.

- ▶ Suppose there exists a drawing of  $K_6$  with two crossings.
- ▶ Both crossings involve four distinct vertices.
- ▶ Since  $K_6$  has six vertices, there is a vertex  $v$  in both crossings.
- ▶ If we delete  $v$ , the resulting graph would have no crossings.
- ▶ This would give a plane drawing of  $K_5$ , a contradiction!

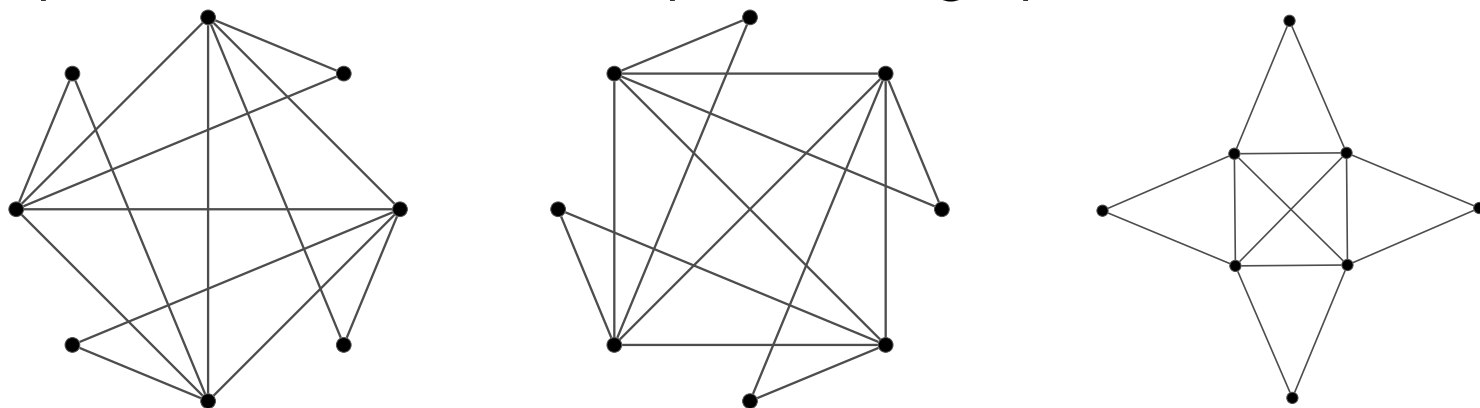
Therefore,  $\text{cr}(K_6) = 3$ .

# The thickness of a graph

**Definition:** The **thickness** of a graph  $G$ , denoted  $\theta(G)$ , is the smallest number of planar subgraphs into which  $G$  can be decomposed. That is, find the optimal way to partition of the edge set of  $G$  into disjoint subsets, each of which is a planar graph.

► So if  $G$  is planar,  $\theta(G) = 1$ , and if  $G$  is non-planar,  $\theta(G) \geq 2$ .

**Example.**  $\theta(K_8) = 2$  since we know  $K_8$  is nonplanar and below is a decomposition of  $K_8$  into two planar subgraphs:



# Theorems about thickness

A simple bound on thickness is:

**Theorem 9.2.1.** If  $G$  has  $p$  vertices and  $q$  edges, then  $\theta(G) \geq \frac{q}{3p-6}$ .

**Proof.** Suppose that  $G = H_1 \cup H_2 \cup \dots \cup H_{\theta(G)}$  is a decomposition of  $G$  into planar subgraphs  $H_i$ , with  $p$  vertices and  $q_i$  edges.

We know that each  $H_i$  must satisfy  $q_i \leq 3p - 6$ . Therefore

$$q = \sum_{i=1}^{\theta(G)} q_i \leq \sum_{i=1}^{\theta(G)} (3p - 6) = \theta(G)(3p - 6).$$

Similarly,

**Theorem 9.2.2.** If  $G$  is a graph with girth  $\geq 4$ , then  $\theta(G) \geq \frac{q}{2p-4}$ .

**Fact:**  $\theta(K_n) = \begin{cases} \left\lfloor \frac{n+7}{6} \right\rfloor & n \neq 9, 10 \\ 3 & n = 9, 10 \end{cases}$  Proved by Beineke,  
Harary, Vasak,  
Alekseev, Gonchakov

# The genus of a graph

A planar graph can always be **embedded on** a sphere.

*That is:* it can be drawn without crossings on the surface of a sphere.

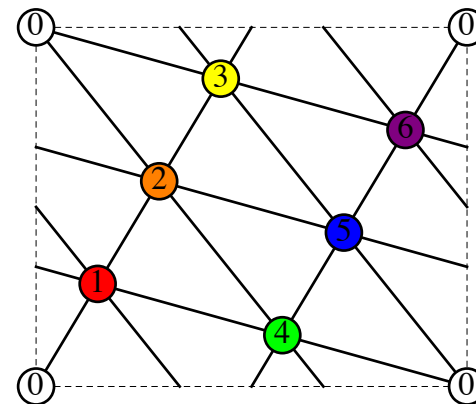
Nonplanar graphs can not be embedded on a plane (or sphere).

What about more complicated surfaces? Like a torus?

**Example.** We can embed  $K_5$  on a torus. (Two ways to see.)

**Example.** We can even embed  $K_7$  on a torus:

However, we can't embed  $K_8$  on a torus. Perhaps on a surface of genus  $g$ ?



# The genus of a graph

*Definition:* The **genus** of a graph is the smallest  $g$  such that  $G$  can be embedded on a surface of genus  $g$  with no crossings.

- ▶ If  $G$  is planar,  $\text{genus}(G) = 0$ ; if  $G$  is non-planar,  $\text{genus}(G) \geq 1$ .

*Fact:* (Ringel, Youngs, 1968) The genus of a complete graph is

$$\text{genus}(K_n) = \left\lceil \frac{(n-3)(n-4)}{12} \right\rceil$$

Embedding on higher genus surfaces changes Euler's formula!

*Theorem.* Let  $G$  be a graph of genus  $g$ . Suppose you have an embedding of  $G$  on a surface of genus  $g$  with no crossings.

If  $r$  is the number of regions, then  $p - q + r = 2 - 2g$ .

*Example.* In our embedding of  $K_5$  on the torus (genus 1):

# Complete graphs

Planarity statistics for complete graphs:

Statistic	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
$cr(K_n)$	0	1	3	9	18	36	60	100	150	225	?	?	?	?	?
$\theta(K_n)$	1	2	2	2	2	3	3	3	3	3	3	3	3	4	4
$genus(K_n)$	0	1	1	1	2	3	4	5	6	8	10	11	13	16	18

The crossing number of a complete graph is unknown for  $n \geq 13$ .

*Conjecture.* (Guy, 1972) The crossing number of a complete graph is

$$cr(G) = \frac{1}{4} \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor \left\lfloor \frac{n-2}{2} \right\rfloor \left\lfloor \frac{n-3}{2} \right\rfloor$$

The cases  $cr(K_{11}) = 100$  and  $cr(K_{12}) = 150$  were proved in **2007**.