

# Stable Marriages

**Let's play matchmaker.**

Given  $n$  men and  $n$  women, where

- ▶ Each man has a complete list of preferences for the women,
- ▶ Each woman has a complete list of preferences for the men,

## Men's Preferences

	Bob	Doug	Fred
1 <sup>st</sup>	Alice	Alice	Elena
2 <sup>nd</sup>	Clara	Elena	Clara
3 <sup>rd</sup>	Elena	Clara	Alice

## Women's Preferences

	Alice	Clara	Elena
1 <sup>st</sup>	Fred	Bob	Doug
2 <sup>nd</sup>	Bob	Fred	Fred
3 <sup>rd</sup>	Doug	Doug	Bob

# Stable Marriages

*Goal:* Create a perfect matching of **stable marriages**.

That is, find a set of  $n$  couples where there are no instabilities:

*Definition:* An **instability** is when one man and one woman both prefer each other to their spouses.

*Example.* Suppose:

- ▶ Alice is married to Bob
- ▶ Clara is married to Doug
- ▶ Bob prefers Clara to Alice

If Clara prefers Bob to Doug: \_\_\_\_\_

If Clara prefers Doug to Bob: \_\_\_\_\_

## The Gale–Shapley Algorithm

*Theorem.* (Gale, Shapley, 1962) In the above situation, there always exists a perfect matching of stable marriages.

*Proof.* Use the **Gale–Shapley Algorithm** to create the marriages.

- 1 Start with no couples engaged.
- 2 As long as at least one man is unengaged, repeat the following:  
Each unengaged man proposes to his next most preferred woman.  
Each woman then decides whether to accept or reject the proposal(s), as follows:
  - ▶ If she is not already engaged, she accepts the proposal.
  - ▶ If she is already engaged, she uses her preference list; she accepts the proposal of the man she prefers the most and rejects all others.
- 3 When all men are engaged, stop. There are  $n$  stable marriages.

<<Time for your moment of zen>>

## Applying the Gale–Shapley Algorithm

Here is a complete set of preferences for 4 men and 4 women.

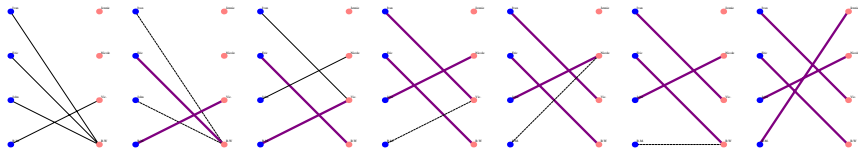
### Men's Preferences

	Ivan	Eric	John	R.M.
1 <sup>st</sup>	R.W.	R.W.	R.W.	Vic.
2 <sup>nd</sup>	Vic.	Jennie	Nicole	Nicole
3 <sup>rd</sup>	Jennie	Vic.	Jennie	R.W.
4 <sup>th</sup>	Nicole	Nicole	Vic.	Jennie

### Women's Preferences

	Jennie	Nicole	Vic.	R.W.
1 <sup>st</sup>	Eric	John	John	Eric
2 <sup>nd</sup>	John	R.M.	Ivan	R.M.
3 <sup>rd</sup>	R.M.	Eric	R.M.	Ivan
4 <sup>th</sup>	Ivan	Ivan	Eric	John

# The Algorithm, Pictorially



## Men's Preferences

Ivan	Eric	John	R.M.
R.W.	R.W.	R.W.	Vic.
Vic.	Jennie	Nicole	Nicole
Jennie	Vic.	Jennie	R.W.
Nicole	Nicole	Vic.	Jennie

## Women's Preferences

Jennie	Nicole	Vic.	R.W.
Eric	John	John	Eric
John	R.M.	Ivan	R.M.
R.M.	Eric	R.M.	Ivan
Ivan	Ivan	Eric	John

## Proof of Correctness

*Claim.* The Gale–Shapley Algorithm gives a set of  $n$  stable marriages.

*Proof.* We must show that the algorithm always stops, and that when it stops, the output is indeed a full set of stable marriages.

### **The algorithm terminates.**

- ▶ In each step, at least one proposal occurs
- ▶ there are a finite number to be made.

*Claim:* Upon termination, everyone is engaged:

- ▶ Once a woman has been proposed to, she stays engaged.
- ▶ If a woman is not engaged at the end, she had no proposal.
- ▶ It follows that there is also some man not engaged; however, he must have proposed to the unengaged woman during some round!

## Proof of Correctness

**The output is a set of stable marriages.**

*We ask:* Is there an instability?

- ▶ Suppose Bob prefers Clara to his current wife.
- ▶ Bob must have proposed to Clara before his current wife.
- ▶ Clara must have turned down Bob.
  - ▶ Clara was proposed to by someone she prefers!
- ▶ Hence, Clara must prefer her current husband to Bob.
- ▶ Therefore, there is no instability.

## Male-optimality

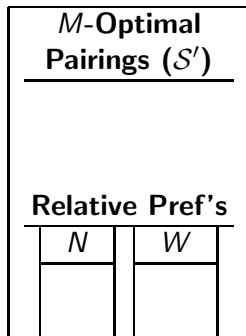
*Claim.* The marriages  $\mathcal{S}$  generated by the Gale–Shapley Algorithm are **male optimal**. That is, given any other set of stable marriages, each man will only be paired with a woman lower on his preference list.

*Proof.* Suppose that during the Gale–Shapley Algorithm, there is a man who is paired with a “sub-optimal” woman.

- Let  $M$  be the first man who is rejected by his optimal woman  $W$  during the algorithm.

[That is, there is some other set  $\mathcal{S}'$  of stable marriages in which  $M$  is paired with  $W$ .]

- $M$  is rejected because some man  $N$  proposes to  $W$  whom  $W$  prefers to  $M$ .
- Since  $M$  is the *first* man rejected, we know  $N$  likes  $W$  at least as much as his optimal woman.
- This, in turn, creates an instability in  $\mathcal{S}'$  since  $W$  prefers  $N$  to  $M$  and  $N$  prefers  $W$  to the woman he is paired with.





## Last remarks

- ▶ The marriages generated by Gale–Shapley are male optimal.
- ▶ The marriages generated by Gale–Shapley are female pessimal.
- ▶ Run the algorithm with the women proposing to reverse the roles.  
If you do this and get the same marriages, \_\_\_\_\_
- ▶ If not all rankings are made, then there may be unmatched people. For example, what if Robot Man did not like Jennie?
- ▶ The National Resident Matching Program (<http://www.nrmp.org>) implements this algorithm to match medical students to residency programs.