Linear Optimization

Today we start our last topic of the semester, linear optimization.

Comprehension goals:

- What is a linear program?
- Visualizing linear programs graphically.
- Understanding solutions graphically.
- Solving linear programs using *Mathematica*

Fertilizer example (p.253)

A fertilizer manufacturer uses nitrates and phosphates to make batches of two different kinds of fertilizer.

- Sod-King fertilizer needs 4 phosphates, 18 nitrates.
- Gro-Turf fertilizer needs 1 phosphate, 15 nitrates.

The profit for one batch of Sod-King is $1000.
The profit for one batch of Gro-Turf is $500.

The company has 10 phosphates and 66 nitrates on hand.

Question. How many batches of each should the company make to earn the most profit?

Initial thoughts?

Fertilizer example (p.253)

Translate the problem into mathematics:
We must determine how many batches to make of each.

- Let x represent the number of batches of Sod-King made.
- Let y represent the number of batches of Gro-Turf made.

What are the constraints on what x and y can be?

- Phosphate constraint:

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What are we trying to maximize?

- Profit:

Linear Programs

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\begin{array}{ll}
\text{Maximize } 1000x + 500y & \\
\text{subject to} & 4x + y \leq 10 \\
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\end{array}
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This is a **linear program**, an optimization problem of the form:

$$\begin{array}{ll} \text{Maximize } c_1x_1 + c_2x_2 + \cdots + c_nx_n \text{ (the \textbf{objective function})} & \\ \text{subject to} & a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n \leq b_1 \\ \text{(the \textbf{constraints}):} & a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n \leq b_2 \\ & \vdots \qquad\qquad \vdots \\ & a_{m1}x_1 + a_{m2}x_2 + \cdots + a_{mn}x_n \leq b_m \end{array}$$

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- The decision variables can have any real value, not only integers.
- All constraints and the objective functions are *linear combinations* of the decision variables. (Coefficients are constants.)
- A linear program in the above form is “easy to solve”.

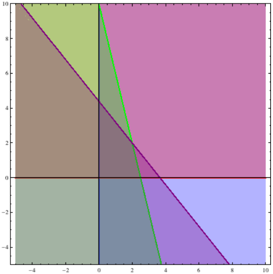
Fertilizer example, graphically

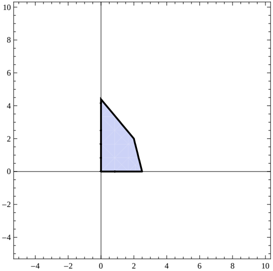
Maximize $1000x + 500y$

subject to the constraints:

$$\begin{aligned} 4x + y &\leq 10 \\ 18x + 15y &\leq 66 \\ x &\geq 0 \\ y &\geq 0 \end{aligned}$$

Let's consider our example graphically.





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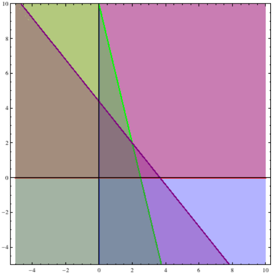

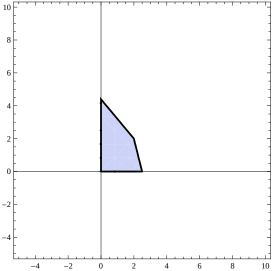

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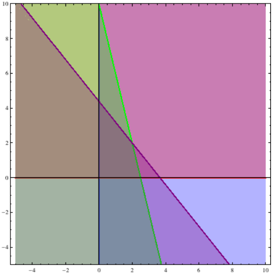

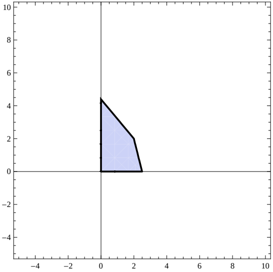

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Definition: The set of points (x, y) that satisfy the constraints is called the **feasible region**.

- In general, points of form $(x_1, x_2, \ldots, x_n)$.
- Feasible region always a polytope. (Always has flat sides and is convex.)
- Feasible region may be bounded or unbounded; might be empty.

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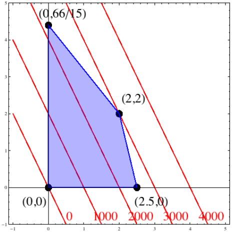

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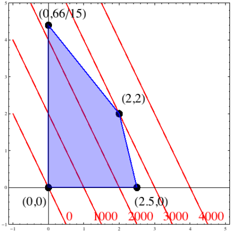

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Is there a point in the feasible region such that $1000x + 500y = 2000$?

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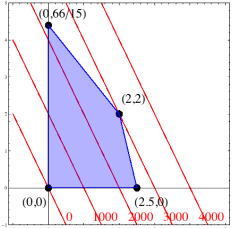

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Is there a point in the feasible region such that $1000x + 500y = 2000$?
Is there a point in the feasible region such that $1000x + 500y = 4000$?

As we plot these **constant-objective** lines, we notice that

- They are parallel.

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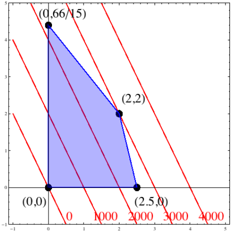

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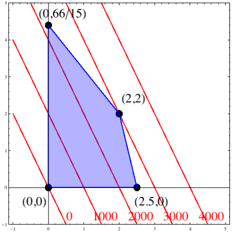

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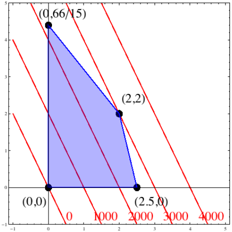

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- They are parallel.
- If there is a feasible region, at least one line will intersect it.
- As we increase the "constant", the last place we touch the feasible region is on the boundary, at one or more corners.

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We have intuited the following theorem.

Theorem. The maximum (or minimum) in a linear program either:

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3. Evaluate the objective function at each corner point.
4. Pick out the optimum value.

Solution of fertilizer example

Maximize $p(x, y) = 1000x + 500y$

subject to the constraints:

$$4x + y \leq 10$$
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$$x \geq 0$$
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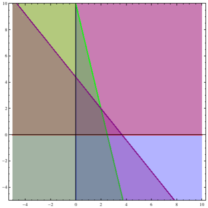

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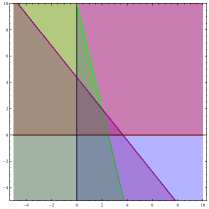

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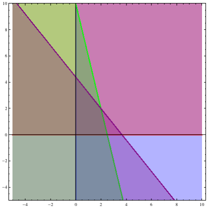

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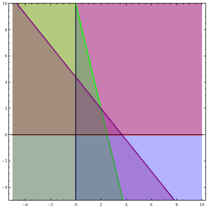

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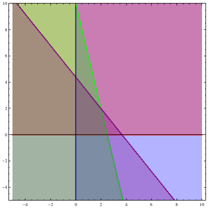

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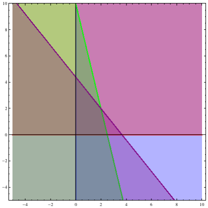

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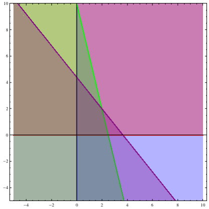

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3. Evaluate the objective function at each corner point.
   - $p(0, 0) = 0$
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   - $p(2, 2) = 3000$.

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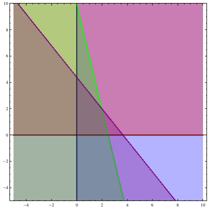

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4. Pick out the optimum value. [Max value: \$3000, occurs at (2,2).]

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```
In[1]: Maximize[{1000 x + 500 y,
         x >= 0 && y >= 0 && 4 x + y <= 10 && 18 x + 15 y <= 66}, {x,y}]
Out[1]: {3000, {x -> 2, y -> 2}}
```

Using *Mathematica* to solve a linear program

Once you have written your optimization problem as a linear program, you can use *Mathematica* to solve your problem.

Use either the `Maximize` or `Minimize` command.

Syntax: `Maximize[{`*obj*`,`*constr*`},`*vars*`]`

- *obj* is the objective function that you wish to optimize.
- *constr* are the set of all constraints, joined with `&&`'s (ANDs).
- *vars* is the set of variables.

```
In[1]: Maximize[{1000 x + 500 y,
         x >= 0 && y >= 0 && 4 x + y <= 10 && 18 x + 15 y <= 66}, {x,y}]
Out[1]: {3000, {x -> 2, y -> 2}}
```

The output gives the optimum value and the values the variables take on there.